

Homotopy

Discrete Morse Theory \leftarrow simplification

Schur Complements

Recall homotopy: $f, g: X \rightarrow Y$

$$f \simeq g \iff \exists h: X \times I \rightarrow Y$$

$$h(\cdot; 0) = f$$

$$h(\cdot; 1) = g$$

In chain category, notion of homotopy:

Chain maps $F_*, G_*: C_* \rightarrow D_*$ ($F_{k-1} \partial_k^c = \partial_k^0 F_k$)

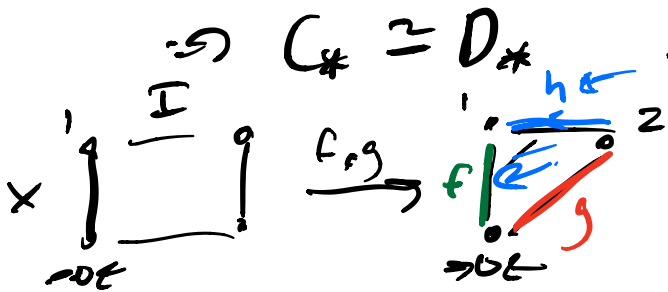
Chain homotopy $H_*: H_k: C_k \rightarrow D_{k+1}$

$$\rightarrow \partial_{k+1}^D H_k + H_{k-1} \partial_k^c = G_k - F_k \Rightarrow F_* \simeq G_*$$

We say C_* and D_* are homotopy equiv.

iff $\exists F_*: C_* \rightarrow D_*$ and $G_*: D_* \rightarrow C_*$

st. $G \circ F \simeq \text{id}^C$, $F \circ G \simeq \text{id}^D$



$$F: \begin{aligned} (0) &\rightarrow (0) \\ (1) &\rightarrow (1) \\ (0,1) &\rightarrow (0,1) \end{aligned}$$

$$G: \begin{aligned} (0) &\rightarrow (0) \\ (1) &\rightarrow (2) \\ (0,1) &\rightarrow (0,2) \end{aligned}$$

$$H: \begin{aligned} (0) &\mapsto (0,0) = 0 \\ (1) &\mapsto (1,2) \leftarrow \leftarrow \\ (0,1) &\mapsto (0,1,2) \end{aligned}$$

in $\text{dim } 0$:

$$\partial H_0 + H_1 \partial = G - F$$

$$(0): 0 + 0 = (0) - (0) = 0 \quad \checkmark$$

$$(1): \partial(1, 2) + 0 = (2) - (1) = G(1) - F(1) \quad \checkmark$$

in $\text{dim } 1$:

$$\partial H_1 + H_0 \partial = G - F?$$

$$-\partial(0, 1, 2) + H_0((1) - (0)) = -(0, 1) - \cancel{(1, 2)} + (0, 2) + \cancel{(1, 2)} + 0$$

$$= -(0, 1) + (0, 2) \\ = -\underbrace{F(0, 1)} + \underbrace{G(0, 1)} \quad \checkmark$$

Prop: if $F_x \cong G_x$, then $H_k(F_x) = \widehat{F}_k \cong H_k(G_x) = \widehat{G}_k$

(homotopic maps induce same map
on homology)

Pf: Let $F_x, G_x: C_x \rightarrow D_x$, $F_x \cong G_x$

let $[x] \in H_k(C_x)$ we want to show

$$[F_k x] = [G_k x]$$

i.e. $G_k x - F_k x \in \text{img } \partial_{k+1}^0$

because $G_x \cong F_x \ni H_x$ s.f.

$$G_k X - F_k X = \underbrace{\partial_{k+1} H_k X}_{\text{e.g. } \partial_{k+1}} + H_k \underbrace{\partial_k X}_{=0 \text{ b/ } X \text{ is } \partial_k \text{ (homology sep)}}$$

✓

\tilde{F}_k, \tilde{G}_k send homology reps to same homology class \rightarrow induced maps are same. \square

Corollary: if $L \simeq D$ then $H_k(L) \cong H_k(D)$

Homotopy invariance of homology:

$$X \simeq Y \Rightarrow C_k(X) \cong C_k(Y) \Rightarrow H_k(X) \cong H_k(Y)$$

$$f \simeq g \Rightarrow F_k \cong G_k \Rightarrow \tilde{F}_k \cong \tilde{G}_k$$

Discrete Morse Theory:

Morse theory connects smooth manifolds w/ cell complexes. Big subject in math.

We'll look at how discrete Morse theory can simplify complexes, ultimately homology calculations.

Ref: Mischaikow, Nanda 2013 ←

Def: Let X be a (simplicial) complex. An incidence function $\kappa: X \times X \rightarrow \mathbb{R}$ pairs simplices w/ codimension-1 faces $\kappa(\sigma, \tau) \neq 0$ if τ is a codim-1 face of σ .

note: δ in chain c_k is an incidence

The face partial order \leq on X is generated by the relation $\sigma \leq \tau$ if $k(\sigma, \tau) \neq 0$

$\rho \leq \sigma$ if $\exists \tau_0 \dots \tau_n$ $\rho \leq \tau_0 < \tau_1 < \dots < \tau_n \leq \sigma$

Def: a partial matching of (X, k) partitions

X into 3 sets: A, K, Q , with bijections

$w: Q \rightarrow K$ s.t. $\forall q \in Q, k(w(q), q) \neq 0$

define $(A, w: Q \rightarrow K)$

$\Rightarrow \dim w(q) = \dim q + 1 \quad q < w(q)$

matching is acyclic b/c w obeys partial order

A gradient path in an acyclic matching:

is a sequence of cells $\rho = (q_1, w(q_1), \dots, q_m, w(q_m))$

with $q_i \in Q$ s.t. $q_i \neq w(q_{i-1}) < w(q_i)$



$(0) < (0,1)$

$(1,2) < (0,1,2)$

$((1,2), (0,1,2), (0), (0,1))$

the index $\nu(\rho)$ is defined

$$\nu(\rho) = \frac{\prod_{i=1}^{m-1} k(w(q_i), q_{i+1})}{\prod_{i=1}^m k(w(q_i), q_i)}$$

initial cell q_i of ρ denoted $q_\rho \in Q$

final cell $w(q_m)$ - - - $k_\rho \in K$

given cells $a, a' \in A$, a gradient path ρ is a connection from a to a' if $q_\rho \prec a$, $a' \prec k_\rho$: $a \hookrightarrow a'$

The multiplicity of the connection ρ is defined $m(\rho) = \kappa(a, q_\rho) \cdot \nu(\rho) \cdot \kappa(k_\rho, a')$

define a new map $\tilde{\kappa} = A \times A \rightarrow \mathbb{R}$ by the relation $\tilde{\kappa}(a, a') = \kappa(a, a') + \sum_{\rho: a \rightarrow a'} m(\rho)$

A complex $(A, \tilde{\kappa})$ is called the Morse complex associated to the acyclic matching $(A, w: Q \rightarrow K)$ of X , and $\tilde{\kappa}$ is called the Morse incidence function. 10 (if no connecting ρ)

Thm: if we consider A as a restriction of X to unmatched cells (simplices) and take $\tilde{\kappa}$ as boundary map of A , then $\underline{H}_*(X) \cong \underline{H}_*(A)$

Example:



$$Q = \{ (0) \}$$

$$K = \{ (0, 1) \}$$

$$A = \{ (1), (2), (1, 2), (0, 2), (0, 1, 2) \}$$

$$w = (0) \mapsto (0, 1)$$

$\varphi = ((0) (0, 1))$ gradient path, $\kappa_\varphi = (0, 1)$

$$q_\varphi = (0)$$

$$\gamma(\varphi) = -1, m(\varphi) = -1$$

ρ : is a connection from (1) and $(0, 2)$

$$\tilde{\delta}((0, 2), (1)) = \delta(\underbrace{(0, 2), (1)}_{(1)}) + \sum_{(0, 2) \xrightarrow{\rho} (1)} m(\varphi)$$

$\rho \leq u$
-1

$$\tilde{\delta}_2 (0, 1, 2) \mapsto (1, 2) - (0, 2)$$

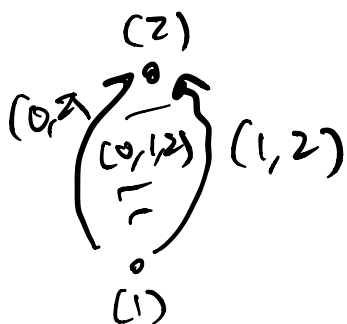
$$\tilde{\delta}_1 : (1, 2) \mapsto (2) - (1)$$

$$(0, 2) \mapsto (2) - (1)$$

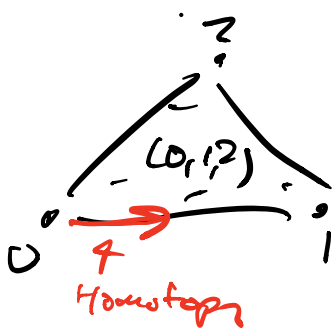
$$\tilde{\delta}_0 : (1) \mapsto 0$$

$$(2) \mapsto 0$$

Corresponding Morse Complex



\cong



Essentially, we performed a retraction of the edge $(0,1)$ to the vertex (1)

PF of thm $H_*(X) \cong H_*(A)$ we proceed by looking at a single retraction at a time
 let $q \in X$, $w: q \mapsto w(q)$ $X_q = X \setminus \{q, w(q)\}$
 $\underbrace{\quad}_A$

$$k_q: X_q \times X_q \rightarrow \mathbb{R}$$

$$k_q(x, y) = k(x, y) - \frac{k(x, q) \cdot k(w(q), y)}{k(w(q), q)}$$

define $F_x: C_*(X) \rightarrow C_*(X_q)$ by

$$F(x) = \begin{cases} 0 & \text{if } x = w(q) \\ - \sum_{y \in X_q} \frac{k(w(q), y)}{k(w(q), q)} y & \text{if } x = q \\ x & \text{otherwise} \end{cases}$$

define $G_x: C_*(X_q) \rightarrow C_*(X)$

$$G(x) = x - \frac{k(x, q)}{k(w(q), q)} w(q)$$

Exercise: both F, G are chain maps.

Lemma: F and G are chain equivalences
if they are maps st. $G \circ F \cong \text{id}$, $F \circ G \cong \text{id}$

PF: $F \circ G = \text{id}^{(x, G)}$ (check)

$G \circ F = \text{id}^{(x, G)}$ chain htpy

$$H_n(x) = \begin{cases} \frac{1}{k(w(q), q)} w(q) & \text{if } x=q \\ 0 & \text{otherwise} \end{cases}$$

This is htpy of q collapsing through $w(q)$.

Can check: $H_n \partial_n + \partial_{n+1} H_n = \text{id} - G \circ F$

Finally, want to check that we can chain
these operations together. i.e. can go through
1 collapse at a time to obtain Morse complex.

Sketch: collapsing a single pair $(q, w(q))$
doesn't affect rest of vector field
and \bar{K}_q are consistent

M + N 2013: extend this to filtrations
using def of filtered acyclic matching
on filtration $\{X^t\}_{t \in T}$

$(A^t, w^t: Q^t \rightarrow K^t)$ for each X^t filtered
 $A^t \subseteq A^{t+1}$, $Q^t \subseteq Q^{t+1}$, $K^t \subseteq K^{t+1}$, $w^t = w^{t+1}|_{Q^t} \neq \in$

this gives a filtration on Morse cpx.

$$M: (A^t, \kappa^t)$$

Furthermore, homological equivalence at every step ensures $PH(X^t) = PH(M^t)$ (up to 0-length)

Schur Complements

Let's revisit the formula for k_q

$$k_q(x, y) = k(x, y) - \frac{k(x, q) k(w(q), y)}{k(w(q), q)}$$

assume that k derived from δ , we'll permute $q, w(q)$ to first row/col respectively.

$$\delta: \begin{array}{|c|} \hline q \\ \hline \end{array} \begin{array}{|c|} \hline w(q) \\ \hline \end{array}$$

recall the Schur complement formula. This shows up in blocked Gaussian elimination.

Let $A = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}$, with A_{11} invertible

$$\text{Then, } A = \begin{array}{c} L \\ \begin{bmatrix} I & 0 \\ A_{21}A_{11}^{-1} & I \end{bmatrix} \end{array} \begin{array}{c} \begin{bmatrix} A_{11} & 0 \\ 0 & S \end{bmatrix} \\ U \\ \begin{bmatrix} I & A_{11}^{-1}A_{12} \\ 0 & I \end{bmatrix} \end{array}$$

S is the "Schur complement"

$$\begin{bmatrix} A_{11} & 0 \\ A_{21} & S \end{bmatrix} \begin{bmatrix} I & A_{11}^{-1} A_{12} \\ 0 & I \end{bmatrix}$$

(1)

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & S + A_{21} A_{11}^{-1} A_{12} \end{bmatrix}$$

$$\Rightarrow S = A_{22} - A_{21} A_{11}^{-1} A_{12}$$

Let's look at a Schur complement of

$$\begin{array}{c} w(q) \\ \begin{array}{|c|} \hline \delta^q \\ \hline \end{array} \end{array} \delta^q [e_{ij}] = \delta [e_{ij}] - \frac{\delta [e_i, w(q)] \cdot \delta [q, z]}{\delta [q, w(q)]}$$

compare to

$$k_q(x, y) = k(x, y) - \frac{k(x, q) k(w(q), y)}{k(w(q), q)}$$

Same form

Lemma from Numerical lin alg

Suppose $A = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}$ with A_{11} invertible.

Then Schur complement $S = A_{22} - A_{21} A_{11}^{-1} A_{12}$

is same whether we use block inverse A_{11}^{-1} or we compute using recursion on invertible submatrices of A_{11}

Schur complement viewpoint on Morse complex
discrete Morse vector field identifies
an invertible block of d_k . Boundary ∂^2
is Schur complement wrt. these blocks.
w/ columns removed (clearing) -
rows removed (compression)
if they are matched in different dimensions.

invertible block identifies basis
for subset of $\text{rang } d_k$

Schur complement just ensures new basis
is independent of this subset.

Final notes:

- Perscus software. demonstrated speedups
- pairs $(q, w(q))$ are persistence pairs.
- speedups can be interpreted as blocking
alg w/ optimizers.
- Ripser & Eivene also use for pre-processing
- Eivene (Heasclman)